



# Ambientes de Desenvolvimento Avançados

<http://www.dei.isep.ipp.pt/~jtavares/ADAV/ADAV.htm>

## Aula 14 Engenharia Informática

2004/2005

José António Tavares  
jrt@isep.ipp.pt

1



## Herança versus Composição

**Capítulo 7 de:**

Szyperski, Clemens et al. Component Software - *Beyond Object-Oriented Programming*. Second Edition

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2



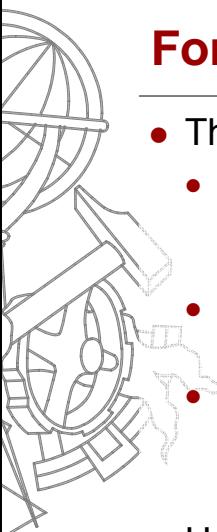
## Conteúdo

- Visão geral
- Vários aspectos da Herança
- Problemas
  - Problemas da classe base frágil
- Abordagens para disciplinar a herança
- Das classes à composição de objectos
  - Reencaminhamento x Delegação

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3



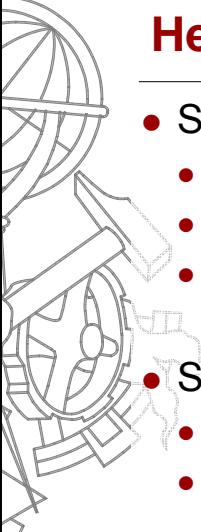
## Formas de Herança

- Three facets of inheritance
  - **Implementation inheritance**  
(sub-classing) sharing of implementation fragments
  - **Interface inheritance**  
(sub-typing) sharing of contract fragments
  - **Substitutivity**  
Promise of substitutability
- How to avoid inheritance ?

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## Herança

- Simula 67 - 1970
  - Inheritance of implementation
  - Inheritance of interfaces
  - Establishment of substitutability
- Smalltalk - 1983
  - Inheritance of implementation
  - Inheritance of interfaces

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## Herança

- Eiffel
  - Possible to undefined inheritance interface feature
- Emerald (1987), Java, C#
  - Interface and implementation inheritance have been separated
- COM and OMG IDL
  - Interface definition language

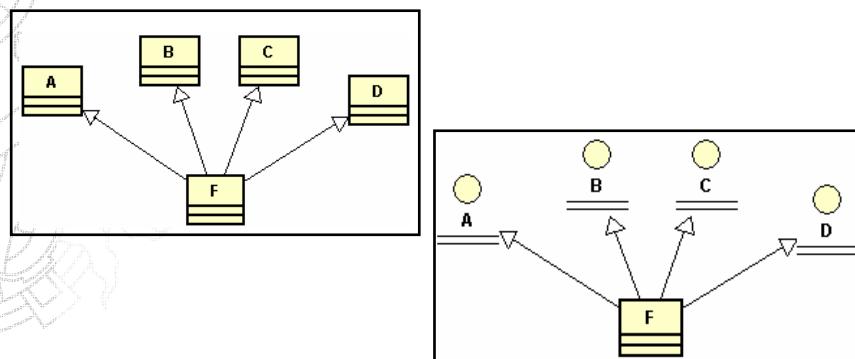
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## More flavors to the soup

- Multiple Inheritance



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## More flavors to the soup

- Multiple Inheritance
  - Establish compatibility with multiple independent context is important. Multiple interface is one way to achieve this.
  - OMG IDL, Java, C# → support multiple interface inheritance
  - COM → not support multiple interface inheritance, but permit that a component support multiple interface simultaneous (that is much the same thing).
  - Multiple interface inheritance does not introduce any major technical problems beyond those already introduced by single interface inheritance.

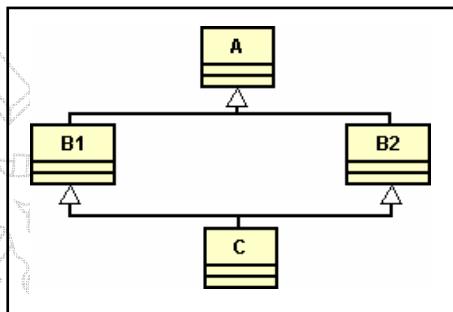
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## More flavors to the soup

- Mixing implementation fragments...



Do both superclasses  
B1 and B2 get their own  
copy of the state  
defined by the  
superclass A?

Diamond inheritance problem

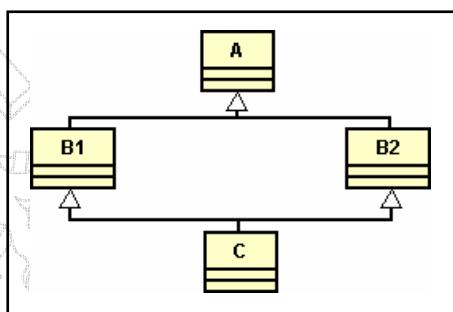
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## More flavors to the soup

- Mixing implementation fragments...



About C class ?

- State...
- Methods ...

Diamond inheritance problem

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## More flavors to the soup

- Some approaches to discipline...
  - CLOS (Common Lisp Object System)
    - Linear order of inheritance
  - C++
    - Maintaining the integrity of sub-objects
  - Java
    - Limited to single implementation inheritance
  - OMG IDL and COM
    - Not support implementation inheritance at all

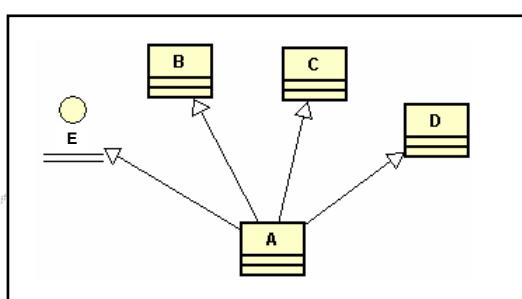
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## More flavors to the soup

- Mixins



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## More flavors to the soup

- Mixins

```
interface B
{
    void X ();
    void Y ();
}
```

```
abstract class X1 implements B
{
    void X () {
        ...
        // X2.Y ();
    }
}
```

```
abstract class X2 implements B
{
    void Y () {
        ...
    }
}
```

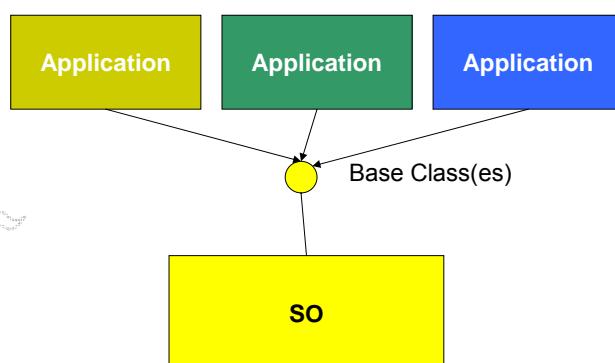
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## Back to basic ingredients...

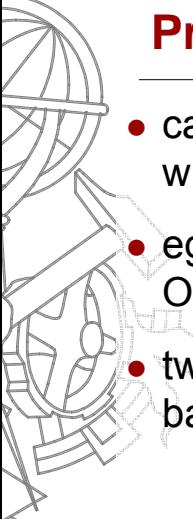
- The Fragile Base Class (FBC) problem



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## Problema da classe base frágil

- can super-class (base class) evolve without breaking subclasses?
- eg old applications with new revision of OS
- two issues: *syntactic* and *semantic* fragile base class problem

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## Problema da classe base frágil

### Syntactic

- a matter of binary compatibility of compiled sub-classes with new binary releases of super-classes
- ‘release-to-release binary compatibility’
- nothing to do with semantics of inherited code
- sub-class should not need recompilation, just because of ‘syntactic’ changes to super-class’s interface
- e.g. moving methods up the class hierarchy
- IBM’s SOM solves this problem by initializing method dispatch tables at load time
- cannot address all ‘syntactic’ changes, e.g. splitting a method in two, or joining two methods into one, or changing a parameter list

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## Problema da classe base frágil

- Semantic

- How can a subclass remain valid in the presence of different version of its super-classes ?
  - Parameters
  - Methods name
  - Return type



Contracts

Versions

Re-entrance

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## Problema da classe base frágil

### Semantic

- how can subclasses remain valid in the face of evolution of the *implementation* of super-classes?
- syntactic FBC addresses problems with immature libraries, but evolution of mature libraries more likely to raise semantic FBC
- to answer this question, it is necessary to understand the semantics of implementation inheritance
- subject of the remainder of this section

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## Up-calls via Herança

- implementation inheritance usually combined with overriding
- selected inherited methods are overridden with new implementations
- new implementations may call overridden code at arbitrary point; abstract methods, or methods of interfaces, may have implementations provided
- invocation of overridden method similar to up-call (method in super-class calling implementation in a sub-class)
- calls span sub-class and super-class in both directions
- but: *every method* is now potentially a callback
- similar problems arise (practical!)
- how to control complexity?

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## Inheritance – more knots than meet the eye

```
abstract class Text
{
    ...
    void write (pos, ch)
    {
        ...
        setCaret(pos);
    }
    void setCaret (int pos)
    {
        caret = pos;
    }
    ...
}
```

```
class SimpleText extends Text
{
    ...
    void setCaret (int pos)
    {
        int old = caretPos();
        if (old != pos)
        {
            hideCaret();
            super.setCaret(pos);
            showCaret();
        }
        ...
    }
}
```

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## Inheritance – more knots than meet the eye

```
abstract class Text
{
    ...
    void write (pos, ch)
    {
        ...
        pos++;
    }
    void setCaret (int pos)
    {
        caret = pos;
    }
    ...
}
```

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```
class SimpleText extends Text
{
    ...
    void setCaret (int pos)
    {
        int oldCaretPos();
        if (oldCaretPos() != pos)
        {
            hideCaret();
            super.setCaret(pos);
            showCaret();
        }
        ...
    }
}
```

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## Abordagens para disciplinar a herança

- these problems known for a while:  
*inheritance breaks encapsulation* (Snyder, 1986)
- early attempts at solution addressed language weaknesses
- but still, sub-class can interfere with and break super-class implementation
- likewise, evolution of super-class can break sub-classes
- some attempts to control use of implementation inheritance:
  - specialization interface
  - partitioning objects
  - reuse contracts (covered in book, not here)

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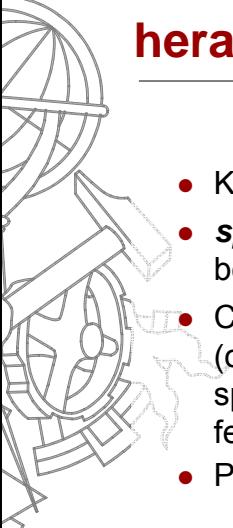
## Abordagens para disciplinar a herança

- The specialization interface

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## Abordagens para disciplinar a herança

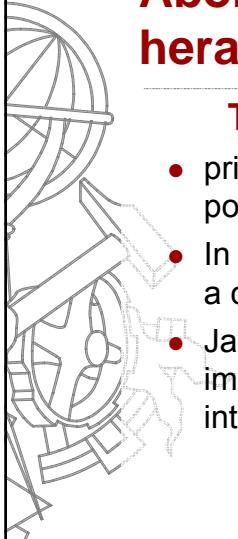
### The specialization interface

- Kiczales and Lamping, 1992
- **specialization interface** is the special interface between class and sub-class
- C++, Java and C#, for example, client interface (outside package) includes only *public* features; specialization interface includes also *protected* features
- Protected - Accessible only to sub-classes

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## Abordagens para disciplinar a herança

### The specialization interface (cont)

- private features can be used to solve problems pointed by Snyder
- In C++, Java and C#, a private feature is private to a class, not an object
- Java, C# and Component Pascal also support the important notion of package-private (or internal) interfaces.

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25



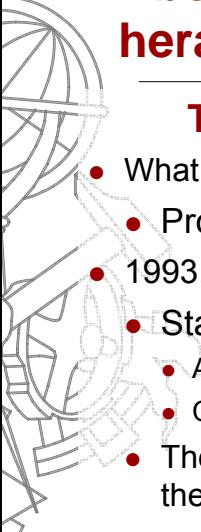
## Abordagens para disciplinar a herança

- Typing the specialization interface

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## Abordagens para disciplinar a herança

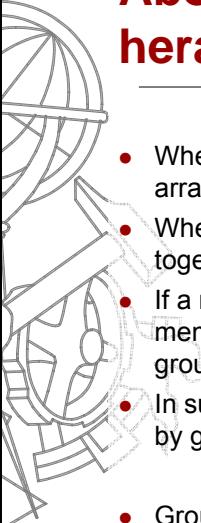
### Typing the specialization interface

- What are the legal modifications a sub-class can apply?
  - Protected interface
  - 1993, John Lamping
  - Statically
    - Acyclic - Arranged in layers
    - Cyclic - Form a group
  - The idea is declare statically which other methods of the same class a given method might DEPEND on.

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## Abordagens para disciplinar a herança

### Typing the specialization interface (cont)

- Where dependencies form acyclic graphs, methods can be arranged in layers;
- Where dependencies form cycles, all the methods in the cycle together form a group;
- If a method need to call another method, it either has to be a member of the called method's group or of a higher layer's group;
- In such an approach, a sub-class has to override methods group by group – either all methods of a group are overridden or none.
- Grouping and layering of methods is seen as a design activity.

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## Abordagens para disciplinar a herança

### Typing the specialization interface (cont)

- The developer determines the **groups or layers**

```
specialization interface Text {  
    state caretRep  
    state textRep  
    abstract posToXCoord  
    abstract posToYCorrd  
    concrete caretPos {caretPos}  
    concrete setCaret {caretRep}  
    concrete write {textRep, caretPos, setCaret}  
    concrete delete {textRep, caretPos, setCaret}  
}
```

No dependencies

- Today no language directly supports Lamping's specialization interface typing

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## Abordagens para disciplinar a herança

- Behavioral specification of the specialization interface

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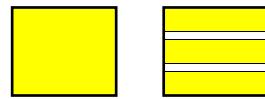
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## Abordagens para disciplinar a herança

### Behavioral specification of the specialization interface

- Lamping's proposal improves information available to sub-classers, but does not address semantic issues of inheritance
- behavioral aspects of inheritance (Stata and Guttag, 1995)
- 1995, Stata & Guttag
  - Class as a combined definition of interacting parts objects
    - Method groups
    - Algebraic specification techniques
    - Notion of behavioral sub-typing



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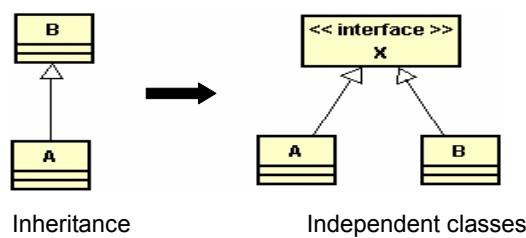
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## Abordagens para disciplinar a herança

### Behavioral specification of the specialization interface (cont)

- to transform ordinary object to Stata-Guttag object group: use only a single sub-object
- Sub-class may change nothing or everything; implementation inheritance useless
- might as well share interface, provide new implementation



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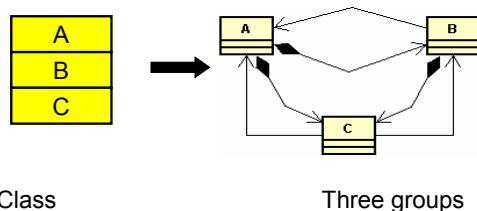
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## Abordagens para disciplinar a herança

### Behavioral specification of the specialization interface (cont)

- conversely, transform object group into collection of objects
- 'self' is lost; how to refer to peers?
- provide each sub-object with references to the others
- to handle object identity, nominate one sub-object the 'main part'



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## Abordagens para disciplinar a herança

- Reuse and cooperation contracts

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## Abordagens para disciplinar a herança

### Reuse and cooperation contracts

- 1996, Steyaert, et. al.
  - Returned to the idea of statically verifiable annotations
  - Reuse contract

Only among methods

```
reuse contract Text {  
    abstract  
        postToXCoord  
        postToYCorrd  
    concrete  
        caretPos  
        setCaret  
        write {caretPos, setCaret}  
        delete {caretPos, setCaret}  
    ...  
}
```

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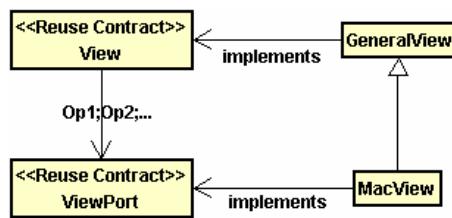
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## Abordagens para disciplinar a herança

### Reuse and cooperation contracts (cont)

- Real innovation - Set of modification operators
  - **Concretization** - replace abstract methods by concrete methods (its inverse is abstraction)
  - **Extension** - add new method that depend on new or existing methods
  - **Refinement** - override methods, introducing new dependencies to possibly new methods.



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## Abordagens para disciplinar a herança

- Representing invariants and method refinements

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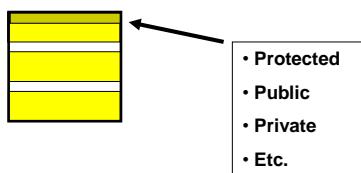
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## Abordagens para disciplinar a herança

### Representation invariants and methods refinements

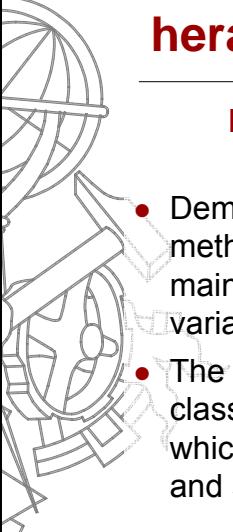
- 1996, Edwards
  - Generalization of the Stata & Guttag
  - Overriding a method in a method group
  - Associating invariants with a class



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38



## Abordagens para disciplinar a herança

### Representation invariants and methods refinements (cont)

- Demonstrate that the overriding of individual methods in a method is permissible if the subclass maintains the representation invariant of the group's variables.
- The idea is to explicitly associate invariants with a class specification that refers to protected variables, which are variables that are only accessible by class and sub-class code (but not external client code)

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39



## Abordagens para disciplinar a herança

- Disciplined inheritance to avoid FBC problems

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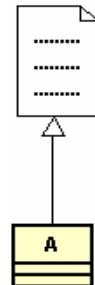
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## Abordagens para disciplinar a herança

### Disciplined inheritance to avoid fragile base class problems

- 1998, Mikhajlov & Sekerinski  
Construir a sub-classe baseada na especificação da super classe, assim a sub-classe ainda será válida mesmo que a implementação da super-classe mude.



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## Abordagens para disciplinar a herança

- Creating correct sub-classes without seeing the super-class code

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## Abordagens para disciplinar a herança

### Creating correct subclasses without seeing superclass code

- 2000, Ruby & Leavens
  - Inverse problem of the semantic FBC problem



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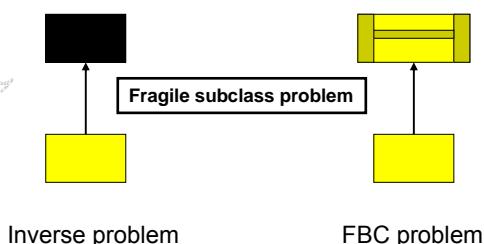
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## Abordagens para disciplinar a herança

### Creating correct subclasses without seeing superclass code (cont)

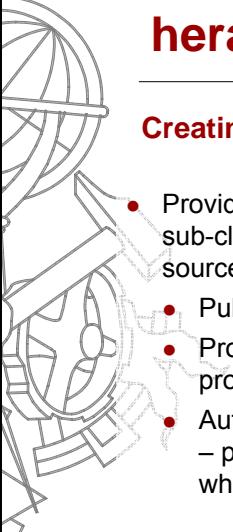
- 2000, Ruby & Leavens
  - Inverse problem of the semantic FBC problem



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## Abordagens para disciplinar a herança

### Creating correct subclasses without seeing superclass code (cont)

- Provide 3 parts to a class specification - for the sub-class can be safely created without requiring access to the source code of the base class
  - Public
  - Protected – reveals information such invariants over protected variables and conditions on protected methods
  - Automatic analysis of the initial source code of the base class – provides information on which variables are accessed and which methods are called by any given method.

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### • Das classes à composição de objectos

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## Das classes à composição de objectos

- Kiczales and Lamping, 1992
- *specialization interface* is the special interface between class and subclass
- eg in Java, client interface (outside package) includes only *public* features; specialization interface includes also *protected* features
- restricts access to interfaces, but doesn't restrict usage by those with access
- distinction between client and descendent interfaces important for controlling implementation inheritance
- sub-class needs to know something about implementation of class

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47

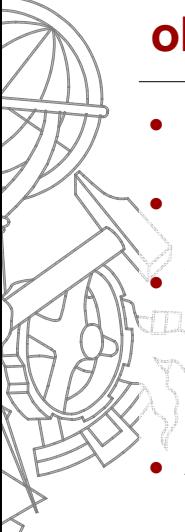
## Das classes à composição de objectos

- motivation for *implementation inheritance* is flexible code reuse
- improving super-class improves sub-classes? re-entrance and up-calls make this difficult
  - *object composition* a simpler alternative ('**has-a**' instead of '**is-a**')
  - **outer object** has the only reference to **inner object**
  - outer object **forwards** messages to inner object
  - improving **inner object** improves **outer object**
- **object composition** and **forwarding** a close approximation to implementation inheritance, without some of the problems

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48



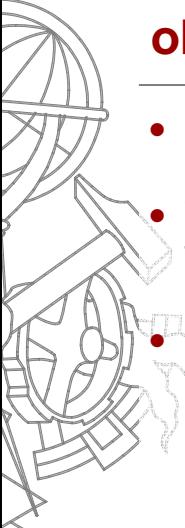
## Das classes à composição de objectos

- Object composition is a much simpler form of composition than implementation inheritance;
- Shares several of the often quoted advantages of implementation inheritance;
- The idea is very simple – whenever an object does not have the means to perform some task locally, it can send messages to other objects, asking for support, and if the helping object is a part of the helped object, this is called *object composition*;
- An object is part of another one if references to it do not leave that object.

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49



## Das classes à composição de objectos

- Sending a message on from one object to another is called **forwarding (re-encaminhamento)**;
- The combination of object composition and forwarding comes fairly close to what is achieved by implementation inheritance;
- However, it does not get so close that it also has the disadvantages of implementation inheritance.

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50

## Das classes à composição de objectos

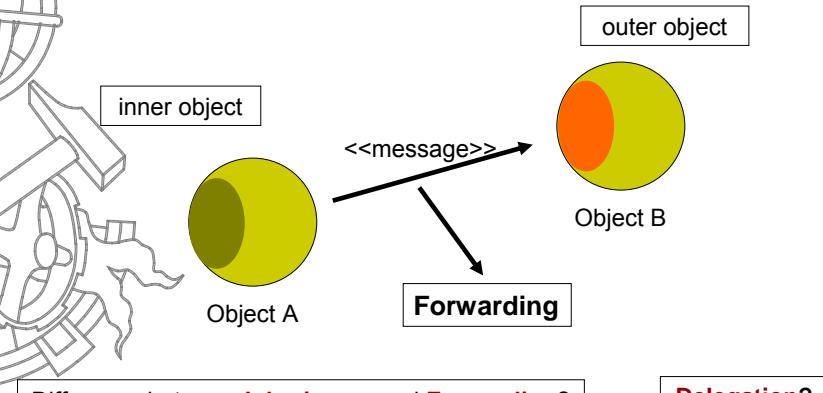
- An **outer object** does not re-implement the functionality of the **inner object** when it forwards messages;
- It reuses the implementation of the inner object;
- If the implementation of the inner object is changed, then this change will “spread” to the outer object;
- The difference between **object composition with forwarding** and **implementation inheritance** is called “**implicit self-recursion**” or “**possession of a common self**”

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51

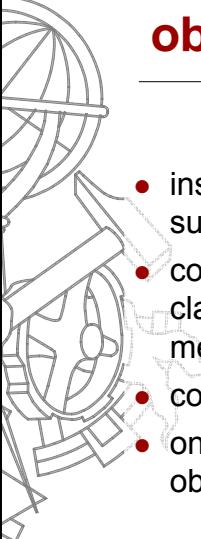
## Das classes à composição de objectos



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## Das classes à composição de objectos

### Possession of a common self

- instance of sub-class shares identity with that of its super-class;
- control can return from a super-class back to a sub-class – invocation of the last overriding version of the method;
- composition of objects has no single identity;
- once control passed from outer to inner object, outer object cannot interfere.

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53



## Das classes à composição de objectos

### Delegation

- **Composition + forwarding** lacks the notion of a common “self”;
- If a common identity is required, it has to be designed in;
- If an object was not designed for composition under a common identity, it cannot be used in such context – mechanisms build in to resend messages to an outer object;
- Object composition supports dynamic and late composition.

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## Das classes à composição de objectos

### Delegation (cont)

- The concept of message passing by delegation is relatively simple;
- Each message-send is classified either as regular send (forwarding) or self-recursive one (delegation)
- Whenever a message is delegated (instead of forwarded), the identity of the first delegator in the current message is remembered;
- Any subsequently delegated message is dispatched back to the original delegator.

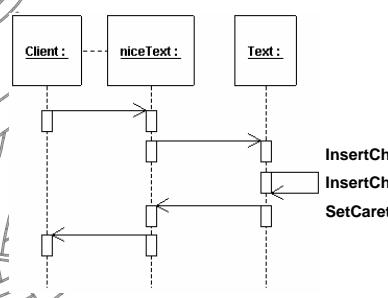
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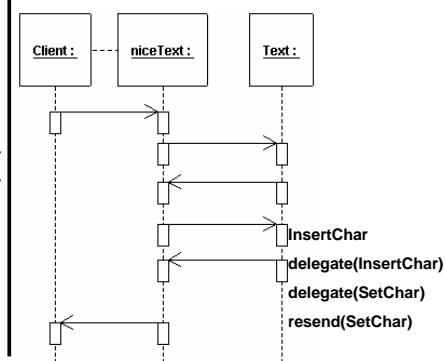
55

## Re-encaminhamento x Delegação

### Forwarding



### Delegation



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56

## Re-encaminhamento x Delegação

### Resumo

- Forwarding
  - Regular Message
- Delegation
  - Self-recursive one
  - Strengthened
  - Identity is remembered
- What the difference between Forwarding and Delegation?



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57

## Delegação x Herança

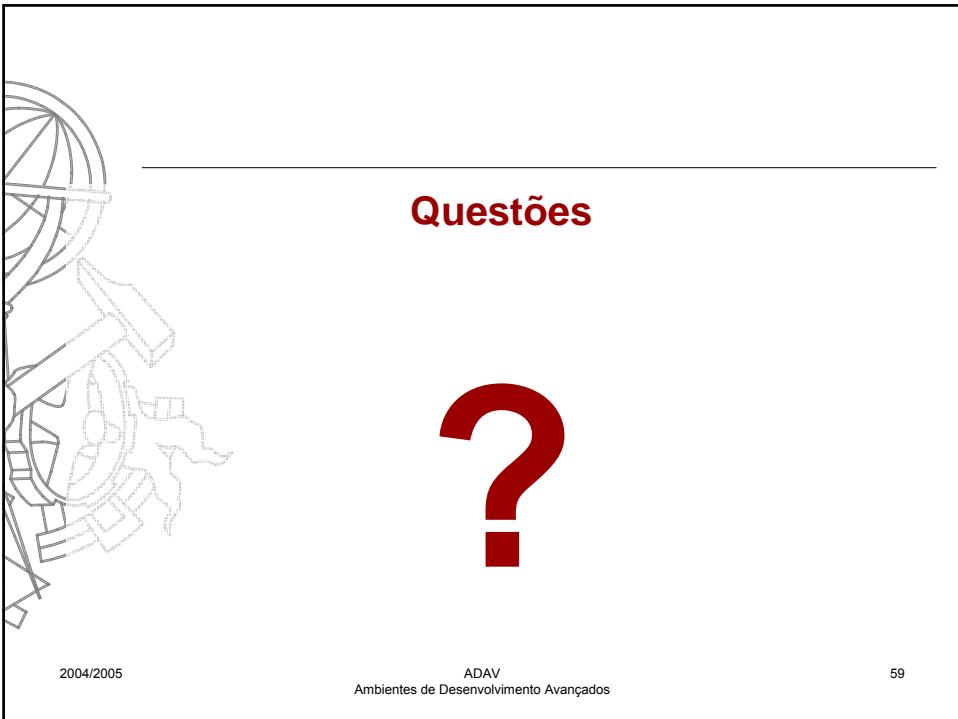
Gamma et al. (1995)

“Delegation has a disadvantage that it shares with other techniques that make software more flexible through object composition: dynamic, highly parameterized software is harder to understand than more static software. [...] Delegation is a good design choice only when it simplifies more than it complicates. [...] Delegation works best when it is used in highly stylized ways – that is, in standard patterns.”

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58



## Questões

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59